



# The Fragmentation Attack in Practice

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Transmit arbitrary WEP data without knowing the key.

- Only requirement: Eavesdrop a single WEP packet.



- 1 Introduction
  - WEP
  - Common Attacks
- 2 Theory
  - PRGA & WEPWedgie
  - Fragmentation
- 3 Practice
  - Hardware & Software Limitations
  - Real-life Attack Example
  - Script-kiddie Tool
- 4 Conclusion

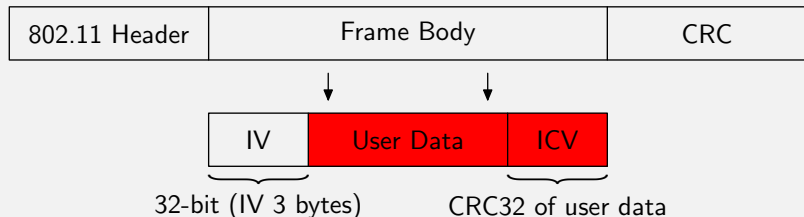
# Wired Equivalent Privacy?

## Overview



- Bogus implementation of RC4 with a 40-bit shared key.
- Only data portion of data packets is encrypted.
- **Initialization Vector** (IV) prepended to key on each encryption.
  - IV is transmitted in clear within WEP packets.

### Data frame format



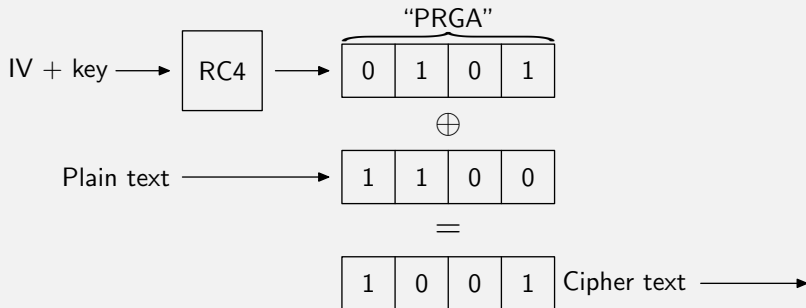
# Wired Equivalent Privacy??

## Encryption



- ① **Seed**: Choose IV (any 24-bit number) and prepend to key.
- ② **KSA**: Run RC4 Key Scheduling Algorithm on seed.
- ③ **PRGA**: Run RC4 Pseudo-Random Generation Algorithm.
- ④ **XOR**: XOR user data with PRGA.

## WEP Encryption





## ① Bruteforce

- 40-bit key!
- ASCII Passphrase.
  - Microsoft Windows XP requires *exactly* 5 or 13 characters.

## ② KSA

- The weak IV attack (aka FMS).
- Requires  $\approx 300,000$ – $3,000,000$  unique IVs.
  - Many networks don't have much traffic.
  - 13% probability IVs improve the attack a lot.
  - aircrack is a good implementation.

## ③ PRGA

- WEP-wedgie: Shared key authentication networks.
- PRGA discovery: Bit-flipping, IV collisions, etc.
- Fragmentation: Not (yet) public!



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### If we had PRGA for an IV:

- Decrypt all packets which use that IV (cipher text  $\oplus$  PRGA).
  - With PRGAs for different IVs, we can decrypt more packets (IV dictionary).
- Encrypt user data with that IV (data  $\oplus$  PRGA).
  - Can always use same IV.

#### Sample PRGA

0	1	0	1	PRGA
0	0	1	1	Plain text
0	1	1	0	Cipher text

If we intercept cipher text and somehow know the clear text:

- Discover PRGA for that IV (cipher text  $\oplus$  clear text).



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- Can decrypt first 128 bytes of packets which use that IV.



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Optimization

Force clients to disconnect by spoofing de-authentication requests—management frames not encrypted!

# PRGA Discovery

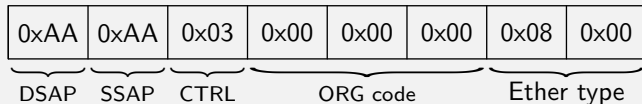
How much clear text do we know?



All data is *Logical Link Control* (LLC) encapsulated.

- Commonly (always) followed by SNAP.
  - Most likely followed by IP.
  - At times followed by ARP.

LLC/SNAP header for IP packet



ARP packets have 0x0806 as ethernet type!

- Distinguishable by fixed and short length.

In general, we can recover at least 8 bytes of PRGA.

# Fragmentation

Greets: Josh Lackey, h1kari, anton, abaddon

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Some details:

- Each fragment needs ICV. Only  $8 - 4 = 4$  bytes for real data.
- Fragment No. field is 4 bits. Only 16 fragments possible.
  - Max data length =  $2^4 \times 4 = 64$ .
  - Can use IP fragmentation too.
- Can generate traffic for which response is known, revealing more PRGA.



- ① Eavesdrop a WEP packet.
- ② Recover 8 bytes of PRGA (clear  $\oplus$  WEP).
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## Speed up other attacks

- ① Send data which generates traffic.
- ② Collect weak IVs.
- ③ Perform KSA attacks (FMS).

## Pure PRGA attack

- ① Send data for which reply is known.
- ② Recover PRGA for more IVs.
- ③ Slowly build an IV dictionary.



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Prism2 (Intersil) based cards.

- Host-AP mode. Can send (almost) raw 802.11 frames.
- Monitor mode. Firmware passes all frames to kernel.

Firmware overwrites 802.11 header fields such as fragment & sequence number!



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Re-write the fields via debug port! (greet to h1kari)

- ① Queue the packet on the card for TX via the normal interface.
- ② Locate the packet on the card's memory via AUX port.
- ③ Instruct the card to begin TX.
- ④ After the firmware processed the header, but before it is sent, overwrite it.
  - In practice, we always win the race!



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- ① Queue and locate packet with 2 random bytes in MAC addr.
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- ③ Overwrite header.

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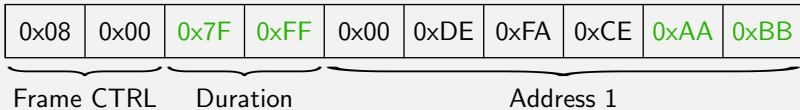


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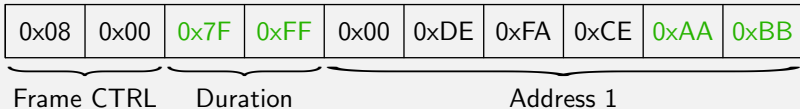


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- Able to send any 802.11 frame and receive all frames.



- **Eavesdrop WEP packet and determine 8 bytes of PRGA.**
- Transmit ARP request (36 bytes) in 9 fragments of 4 data bytes.
  - Who has 192.168.0.1 tell 192.168.0.123.
- Didn't get any reply.
  - Wrong IP network.
  - But AP relayed the packet (since it's a broadcast).
  - Re-encrypted by the AP.
  - Knowing the contents, we discover 36 bytes of PRGA.
- Send ARP request padded with  $x$  0s (in larger fragments).
  - AP relays the longer ARP request.
  - Discover  $36 + x$  bytes of PRGA.
  - Repeat until, say, 1504 bytes of PRGA are known.
- Can send 1500 bytes of data *without* fragmenting.



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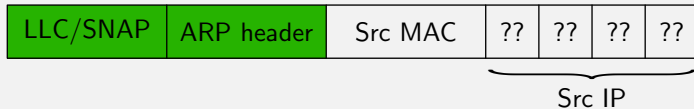
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### ARP decryption

Know whether its ARP request/reply depending on whether its a broadcast or not.



















## By sending ARP request for 192.168.11.1

- Know MAC of router (clear in 802.11 header).
- Router knows our MAC/IP pair (ARP backward learning).

Send ICMP echo to a host we own on Internet.

- Use “our” source MAC/IP pair.
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### Generate traffic to speed up KSA attacks

- Cause controlled host on Internet to flood network.
- Send ARP requests and ICMPs to broadcast IP.
  - Could generate  $\approx 200$  packets/s of traffic.
- Key was actually 40-bit alpha-numeric ASCII.
  - Bruteforceable in  $\leq 5$  minutes ...

### Login to AP and clean up

- Default passwords work great. (root without password here.)
- Clear the logs.
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Designed for Atheros based cards.

- Queue the packet and it shall be sent—No firmware hacks!
- Supports 802.11 a/b/g.
- FreeBSD ath driver patched to support injection.
  - Problem with sending 802.11 ACKs. Possibly they are sent too late—DIFS rather than SIFS.
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IV dictionary built in parallel!

Binds to a TAP interface allowing transmission and reception (if PRGA is known).



After a single ARP request is eavesdropped:

- 144 bytes of PRGA are recovered in 1 second.
- IP is decrypted in  $< 30$  seconds.
- Internet host is contacted in  $< 1$  minute (total time).



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### Traffic generation rate

Flood source	$\approx$ p/s
802.11b client FTP download.	150
LAN client ping -f (no replies).	550
Internet flood (MTU sized packets). ARP replay.	250 350
Internet flood (short packets).	950

Full dictionary requires  $\approx \frac{2^{24}}{250} \times \frac{1}{3600} \approx 18.6$  hours of flooding.



# The Tool: wesside

Key recovery time

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## Total attack time for /dev/urandom keys

Key	Packets	Time (m)
2C:CE:FC:1D:2B	100,000	1.93
80:19:B8:3F:C8	200,000	3.83
6F:34:11:BC:A3	200,000	4.30
91:B7:C0:A7:F7	300,000	5.45
3B:07:DA:02:B7	300,000	5.60
EB:A6:50:D0:2B:DA:CC:B7:E1:B7:E8:50:59	1,700,000	30.77
D9:06:CA:9E:EA:B3:18:CD:24:9F:2E:5E:10	2,400,000	42.85
5E:02:F4:83:FE:F6:27:10:21:EC:8E:87:27	2,700,000	49.17
64:AC:EE:55:B7:7E:27:93:09:6B:78:00:78	9,000,000	156.58
41:0A:68:52:5B:BE:C7:64:D7:09:FC:CC:BB	10,000,000	181.28

# The Tool: wesside

Screen shot

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```
# ./wesside -s 1.2.3.4
[10:49:50] Setting up ath0... done
[10:49:50] Opened tap device: tap3
[10:49:50] Set tap MAC to: 00:00:DE:FA:CE:0D
[10:49:50] Looking for a victim...
[10:49:53] Found SSID(sorbo) BSS=(00:06:25:FF:D2:29) chan=11
[10:49:53] Authenticated
[10:49:53] Associated (ID=3)
...
```

# The Tool: wesside

Screen shot

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...

[10:49:54] Got ARP request from (08:00:46:9E:AF:CD)

[10:49:54] Got 8 bytes of prga IV=(42:bc:00)

[10:49:54] Got 36 bytes of prga IV=(43:bc:00)

[10:49:55] Got 144 bytes of prga IV=(52:bc:00)

[10:49:58] Guessing PRGA 5f (IP byte=255)

[10:49:58] Got clear-text byte: 192

[10:50:00] Guessing PRGA 2d (IP byte=175)

[10:50:00] Got clear-text byte: 168

[10:50:09] Guessing PRGA f7 (IP byte=0)

[10:50:09] Got clear-text byte: 1

[10:50:18] Guessing PRGA f7 (IP byte=102)

[10:50:18] Got clear-text byte: 100

[10:50:18] Got IP=(192.168.1.100)

[10:50:18] My IP=(192.168.1.123)

[10:50:18] Sending arp request for: 192.168.1.1

[10:50:18] Got arp reply from (00:06:25:FF:D2:27)

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[10:49:58] Guessing PRGA 5f (IP byte=255)
[10:49:58] Got clear-text byte: 192
[10:50:00] Guessing PRGA 2d (IP byte=175)
[10:50:00] Got clear-text byte: 168
[10:50:09] Guessing PRGA f7 (IP byte=0)
[10:50:09] Got clear-text byte: 1
[10:50:18] Guessing PRGA f7 (IP byte=102)
[10:50:18] Got clear-text byte: 100
[10:50:18] Got IP=(192.168.1.100)
[10:50:18] My IP=(192.168.1.123)
[10:50:18] Sending arp request for: 192.168.1.1
[10:50:18] Got arp reply from (00:06:25:FF:D2:27)
...
```

# The Tool: wesside

Screen shot

Introduction  
Theory  
Practice  
Conclusion



...

[10:51:28] WEP=000100460 (next crack at 100000) (rate=1448)

[10:51:28] Starting crack PID=17410

[10:52:28] WEP=000185271 (next crack at 200000) (rate=1426)

[10:52:28] Stopping crack PID=17410

[10:52:39] WEP=000201124 (next crack at 200000) (rate=1433)

[10:52:39] Starting crack PID=17412

[10:52:40] WEP=000203778 (next crack at 300000) (rate=1365)

[10:52:41] KEY=(2C:CE:FC:1D:2B)

Owned in 2.85 minutes

#



- Able to transmit arbitrary data on most (all?) 802.11 WEP networks after having eavesdropped a single data packet.
- Can potentially recover a WEP key in a couple of hours.

## Future Work:

- Develop method for higher flood rates (p/s).
- Study how IV generator can be reset—smaller dictionaries.
- Implement a more sophisticated tool and make a Live CD!

A final thought for the adventurous. . .

Assume the AP uses default password for WWW interface.

- Connect to WWW and request WEP configuration page.
- Decrypt TCP sequence number for connection ACK.
- Decrypt contents of page returned—may contain WEP key!